

PwA Cheatsheet

Common Distributions

Discrete				
Name	pmf	cdf	mean	variance
Binomial(n,p)	$\binom{n}{k} p^k (1-p)^{n-k}$	$F(k; n, p) = \Pr(X \leq k) = \sum_{i=0}^{\lfloor k \rfloor} \binom{n}{i} p^i (1-p)^{n-i}$	np	$np(1-p)$
Neg. Binomial(r,p)	$\binom{i-1}{r-1} p^r (1-p)^{i-r}$	-	$\frac{r}{p}$	$r \frac{1-p}{p^2}$
Bernoulli(p)	$\begin{cases} q = (1-p) & \text{for } k = 0 \\ p & \text{for } k = 1 \end{cases}$	$\begin{cases} 0 & \text{for } k < 0 \\ 1-p & \text{for } 0 \leq k < 1 \\ 1 & \text{for } k \geq 1 \end{cases}$	p	$p(1-p)$
Uniform(a,b)	$\frac{1}{n}, n = b - a + 1$	$\frac{ k - a + 1}{n}$	$\frac{a+b}{2}$	$\frac{(b-a+1)^2 - 1}{12}$
Geometric(p)	$p(1-p)^{i-1}$	$1 - (1-p)^i$	$\frac{1}{p}$	$\frac{1-p}{p^2}$
Hypergeometric(N,K,n) "k successes in N, K successful in N"	$\frac{\binom{K}{k} \binom{N-K}{n-k}}{\binom{N}{n}}$	-	$n \frac{K}{N}$	$n \frac{K}{N} \frac{(N-K)}{N} \frac{N-n}{N-1}$
Poisson(λ)	$\frac{\lambda^k e^{-\lambda}}{k!}$	$e^{-\lambda} \sum_{i=0}^{\lfloor k \rfloor} \frac{\lambda^i}{i!}$	λ	λ

Events	
Sample Space	$S = \{\text{all possible outcomes}\}$
Event	$E \subset S$
Union (either or both)	$E \cup F$
Intersection (both)	$E \cap F$ or EF
Complement	$E^C = S \setminus E \Rightarrow P(E^C) = 1 - P(E)$
Inclusion-Exclusion	$\hookrightarrow P(A \cup B) = P(A) + P(B) - P(A \cap B)$
DeMorgan's Law	<ol style="list-style-type: none"> 1. $(E_1 \cup \dots \cup E_n)^C = E_1^C \cap \dots \cap E_n^C$ 2. $(E_1 \cap \dots \cap E_n)^C = E_1^C \cup \dots \cup E_n^C$ <ol style="list-style-type: none"> 1. $0 \leq P(E) \leq 1$ 2. $P(S) = 1$ 3. For mutually excl. events $A_i, i \geq 1$: $P(\bigcup_{i=1}^{\infty} A_i) = \sum_{i=1}^{\infty} P(A_i)$
Finite S, Equal Probability for all point sets:	$P(A) = A / S $
Odds of Event	$\alpha = \frac{P(A)}{P(A^C)} = \frac{P(A)}{1-P(A)}$

Conditional Probability and Independence I	
Conditional Probability	$P(F E) = \frac{P(F \cap E)}{P(E)}$
Independence if	$P(F \cap E) = P(F)P(E)$
Multiplication Rule	$P(E_1 E_2 \dots E_n) = P(E_1)P(E_2 E_1) \dots P(E_n E_1 \dots E_{n-1})$
Bayes Formula (simple)	$P(A B) = \frac{P(B A)P(A)}{P(B)}$
Bayes Formula (full)	$P(A_i B) = \frac{P(B A_i)P(A_i)}{\sum_j P(B A_j)P(A_j)}$
Conditional pmf (discrete)	$p_{X Y}(x y) = \frac{p(x,y)}{p_Y(y)}$
Conditional pdf (discrete)	$f_{X Y}(x y) = \sum_{a \leq x} p_{X Y}(a y)$
Conditional Density (continuous)	$f_{X Y}(x y) = \frac{f(x,y)}{f_Y(y)}$
Conditional Probabilities (continuous)	$P\{X \in A Y = y\} = \int_A f_{X Y}(x y) dx$

Random Variables (Discrete)	
Distribution Function	$F(x) = P\{X \leq x\}$
Probability Mass Function	$p(x) = P(X = x)$
Joint Probability Mass Function	$P(X = x \text{ and } Y = y) = P(Y = y X = x) \cdot P(X = x)$ $= P(X = x Y = y) \cdot P(Y = y)$
Expectation	$E[X] = \sum_{x:p(x)>0} x p(x)$ $\hookrightarrow \text{note: } E[g(X)] = \sum_{x:p(x)>0} g(x)p(x)$
Variance	$Var(X) = E[(X - E[X])^2]$ $= E[X^2] - (E[X])^2$
Standard Derivation	$\sigma = \sqrt{Var(X)}$
Covariance	$Cov(X, Y) = E[(X - E[X])(Y - E[Y])]$ $= E[XY] - E[X]E[Y]$
Moment Gen. Function	$M(t) = E[e^{tX}]$ (same for continuous RVs)

Random Variables (Continuous) I	
Probability Density Function	f such that $P\{X \in B\} = \int_B f(x)dx$
Distribution Function	F such that $\frac{d}{dx}F(x) = f(x)$
Expectation	$E[X] = \int_{-\infty}^{\infty} xf(x)dx$ $\hookrightarrow \text{note: } E[g(X)] = \int_{-\infty}^{\infty} g(x)f(x)dx$
Variance	$Var(X) = E[(X - E[X])^2]$ $= E[X^2] - (E[X])^2$
Standard Derivation	$\sigma = \sqrt{Var(X)}$
Covariance	$Cov(X, Y) = E[(X - E[X])(Y - E[Y])]$ $= E[XY] - E[X]E[Y]$

Random Variables (Continuous) II	
Marginal pmfs	$f_X(x) = \int_{-\infty}^{\infty} f(x, y)dy$ $f_Y(y) = \int_{-\infty}^{\infty} f(x, y)dx$
More on Expectation, Variance, ..	$E[X + Y] = E[X] + E[Y]$ $E[\alpha X] = \alpha E[X]$ $Var(X + a) = Var(X)$ $Var(aX + b) = a^2 Var(X)$ $Var(X + Y) = E[(X + Y)^2] - (E[X + Y])^2$ $= E[X^2 + 2XY + Y^2] - (E[X] + E[Y])^2$ $= E[X^2] + 2E[XY] + E[Y^2] - (E[X])^2 - 2E[X]E[Y] - (E[Y])^2$ $= Var(X) + Var(Y) + 2(E[XY] - E[X]E[Y])$ $= Var(X) + Var(Y) + 2(Cov(X, Y))$ $\Rightarrow E[f(X)g(Y)] = E[f(X)]E[g(Y)]$ $\Rightarrow E[XY] = E[X]E[Y]$ $\Rightarrow Cov(X, Y) = 0$ $\Rightarrow Var(X + Y) = Var(X) + Var(Y)$

Correlation $corr(X, Y) = \rho(X, Y) = \frac{Cov(X, Y)}{\sqrt{Var(X)Var(Y)}}$	<ol style="list-style-type: none"> 1. $-1 \leq \rho(X, Y) \leq 1$ 2. Independence $\Rightarrow \rho(X, Y) = 0$ $Y = mX + cm, m \neq 0$ and c: 3. $m > 0 \Rightarrow \rho(X, Y) = 1$ $m < 0 \Rightarrow \rho(X, Y) = -1$
$E[X] = E[E[X Y]]$	
Disc.: $E[X] = \sum_y E[X Y = y]P\{Y = y\}$	
Cont.: $E[X] = \int_{-\infty}^{\infty} E[X Y = y]f_Y(y)dy$	

Combinatorial Analysis

Order matters and k = n	Permutation
Order does matter and k < n	Variation
Order does not matter and k < n	Combination

Counting

Basic Counting Principle	Experiments E_1, E_2, \dots, E_r with n_1, n_2, \dots, n_r possible outcomes. Total outcomes: $\prod_i^n n_i$
Permutations (without Repeats)	$n! = n \cdot (n-1) \cdot \dots \cdot 1$
Permutations (with Repeats)	$\frac{n!}{k!} = n \cdot (n-1) \cdot \dots \cdot (k+1)$
Variations (without Repeats)	$n \cdot (n-1) \cdot \dots \cdot (n-k+1) = \frac{n!}{(n-k)!}$
Variations (with Repeats)	$\underbrace{n \cdot \dots \cdot n}_{k\text{-times}} = n^k$
Combinations (without Repeats)	$\frac{n!}{(n-k)!k!} = \frac{n(n-1)(n-2)\dots(n-k+1)}{k!} = \binom{n}{n-k} = \binom{n}{k}$ "Binomial Coefficient"
Multinomial Coefficient	$\frac{n!}{n_1!n_2!\dots n_r!} = \binom{n!}{n_1, n_2, \dots, n_r}$ "divide n into r non-overlapping subgroups of sizes n_1, n_2, \dots "
Combinations (with Repeats)	$\frac{(n+k-1)!}{(n-1)!k!} = \binom{n+k-1}{k} = \binom{n+k-1}{n-1}$

Limit Theorems

$$Z_n = \frac{((X_1 + X_2 + \dots + X_n) - n\mu)}{\sigma\sqrt{n}}$$

Central Limit Theorem

$$\text{Then as } n \rightarrow \infty \quad P(Z_n \leq x) \rightarrow \int_{-\infty}^x \frac{1}{\sqrt{2\pi}} \exp(-\frac{1}{2}u^2) du$$

i.e. $P(Z_n \leq x) \rightarrow P(Y \leq X)$ where $Y \sim N(0, 1)$
 $E[X_i] = \mu \quad \text{Var}(X_i) = \sigma^2$

Weak Law of Large Numbers

$$s_n = \frac{1}{n}(X_1 + \dots + X_n)$$

then for any $\epsilon > 0$

$$\lim_{n \rightarrow \infty} P(|s_n - \mu| > \epsilon) = 0$$

Strong Law of Large Numbers
Markov's Inequality

$$P\{\lim_{n \rightarrow \infty} (X_1 + X_2 + \dots + X_n) / n = \mu\} = 1$$

$$P\{x \geq a\} \leq \frac{E[X]}{a}$$

Chebyshev's Inequality

$$E[Y^2] < \infty, \forall a > 0.$$

$$P(|Y| \geq \frac{1}{a^2} E[Y^2])$$

One-sided Cheby-shev (mean 0)

$$P\{|X - \mu| \geq a\} \leq \frac{\sigma^2}{a^2}$$

$$P\{X \geq a\} \leq \frac{\sigma^2}{\sigma^2 + a^2}$$

$$P\{X \geq a\} \leq e^{-ta} M(t) \quad t > 0$$

Chernoff Bounds

$$P\{X \leq a\} \leq e^{-ta} M(t) \quad t < 0$$

Markov Chains

Discrete

$P_{i,j} = P(\text{system is in state } j \text{ at time } n+1 \mid \text{system in state } i \text{ at time } n)$

$$\text{Transition Matrix: } P = \begin{pmatrix} p_{1,1} & p_{1,2} & \dots & p_{1,n} \\ p_{2,1} & p_{2,2} & \dots & p_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ p_{n,1} & p_{n,2} & \dots & p_{n,n} \end{pmatrix}$$

Probability vector $\pi^{(n)}$: Probabilities that we are in state i at n.

$$\pi^{(n+1)} = \pi^{(n)}P$$

$$\pi^{(n)} = \pi^{(0)}P^n$$

A Markov Chain is ergodic (aperiodic and irreducible) iff there exists

$n \in \mathbb{N}^+$ such that P^n has no zero entries. It then has a Steady State

Probability vector $\pi = \lim_{n \rightarrow \infty} \pi^{(n)}$ independent of $\pi^{(0)}$.

1. $\pi_0 + \pi_1 + \dots + \pi_N - 1 = 1$
2. $\pi = \pi P$

Continuous

Poisson

$$P(\tilde{N}(t) = k) = \frac{(\lambda t)^k}{k!} e^{-\lambda t} \text{ if}$$

1. For any fixed t , $\tilde{N}(t)$ is a discrete RV
2. $\tilde{N}(0) = 0$
3. # of events in disjoint intervals are independent
4. $\tilde{N}(t+h) - \tilde{N}(t) = \# \text{ of events in } [t, t+h]$ for $h \rightarrow 0$
5. $P(\tilde{N}(h) = 1) = P(\text{event occurs in } [t, t+h]) = \lambda h + E(h)$
 $(E(h)/h \rightarrow 0, \text{ as } h \rightarrow 0)$
6. $\frac{1}{h} P(\tilde{N}(h) \geq 2) \rightarrow 0, \text{ as } h \rightarrow 0$

Birth-Death

Birth Rates $\lambda_{i,i+1} = b_i$

Death Rates $\lambda_{i,i-1} = d_i$

$\lambda_{i,j} = 0$, otherwise

→ Have steady state prob. vector if b_i s and d_i s are non-zero and we have a finite number of states.

$$1. \pi_0 + \pi_1 + \dots + \pi_N - 1 = 1$$

$$2. \pi_j = \frac{b_0 \dots b_{j-1}}{d_0 \dots d_{j-1}} \pi_0$$

M/M/S Queue

Customers arrive with Poisson Process rate λ , S servers, Service time exponentially distributed with mean $\frac{1}{\mu}$. State j = j customers in queue,

$$b_j = \lambda, d_j = \begin{cases} j\mu, & j = 1, 2, \dots, S \\ S\mu, & j \geq S \end{cases}$$

→ Has steady state prob. vector if $\lambda < S\mu$.

M/M/1 Queue

$$\pi_j = (1 - \frac{\lambda}{\mu})(\frac{\lambda}{\mu})^j, \text{ Mean Queue Length } E[J] = \frac{\lambda}{\mu - \lambda}$$

Surprise, Uncertainty & Entropy

$$H(X) := -\sum_k p_x(x_k) \log_2 p_x(x_k)$$

Entropy

$$(0 \log_2(0)) := 0$$

Surprise

$$S(X = x_k) = -\log_2 p_x(x_k)$$

1. $S(1) = 0 \neq S(0)$ (which is undefined)

2. S decreases: $p < q \Rightarrow S(q) < S(p)$

3. $S(pq) = S(p) + S(q)$

↪ Properties If S is continuous and these are satisfied, $\exists c > 0. \forall p \in [0, 1], S(p) = -c \log(p)$

Average

Uncertainty

$$H(X, Y) := -\sum_j \sum_k p_{X,Y}(x_j, y_k) \log_2 p_{X,Y}(x_j, y_k)$$

Uncertainty of X given Y

$$H_{Y=y_k}(X) := -\sum_j p_{X|Y=y_k}(x_j) \log_2 p_{X|Y=y_k}(x_j)$$

Conditional

Entropy

$$H_Y(X) := \sum_k H_{Y=y_k}(X)p_Y(y_k)$$

Coding Theory

Code C

A map from $\{x_k\} \subset \mathbb{R}$ into sequences of 0's and 1's. Sequences are called code words.

Code Word length

$$x_k \mapsto 0111 \Rightarrow n + k =$$

Expected Length of Code C

$$E[\mathcal{C}] = \sum_k n_k p_k = \sum_k n_k P(X = x_k)$$

No code word extends another one:

✗	✓
$x_1 \mapsto 0$	$x_1 \mapsto 0$
$x_2 \mapsto 00$	$x_2 \mapsto 10$

For any acceptable code assigning n_k bits to x_k the following holds:

Noiseless Coding Theorem

$$E[\mathcal{C}] = \sum_k n_k p_k \geq H(X) = -\sum_k p_k \log_2 p_k$$

Where $p_k = P(X = x_k)$ and n_k length of a codeword associated with x_k

For any discrete RV X there exists an acceptable code with the expected length $E[\mathcal{C}] = L$ such that

↪ Thrm

$$H(X) \leq L < H(X) + 1$$

Algorithm for finding an acceptable code with expected length

$$H(X) \leq E[\mathcal{C}] = L < H(X) + 1 \text{ for discrete RV X:}$$

1. Let n_j be the integer satisfying $-\log_2 p_j \geq n_j < -\log_2 p_j + 1$
2. Find any acceptable code assigning n_j bits to x_j

There is no unique nearly-optimal code in general. Optimal or nearly-optimal coding depends on the pmf of X.

Common Moment Generating Functions M(t)

Binomial

$$(pe^t + 1 - p)^n$$

Neg. Binomial

$$[(pe^t) \div (1 - (1-p)e^t)]^r$$

Poisson

$$\exp(\lambda(e^t - 1))$$

Uniform

$$(e^{tb} - e^{ta}) \div (t(b-a))$$

Exponential

$$\lambda \div (\lambda - t)$$

Normal

$$\exp(\mu t + ((\sigma^2 t^2) \div 2))$$